

November 22nd, 2013 Córdoba, Argentina FaMAF

# **BEYOND HEISENBERG: Entropic and Fidelity-Based Uncertainty Relations** Gustavo Martin BOSYK IFLP & Dto. de Física, UNLP





# THE UNCERTAINTY PRINCIPLE



Heisenberg, 1927: Given two (or more) quantum observables one cannot predict with certainty and simultaneously the outcomes of both

### **UNCERTAINTY RELATION: GENERAL FORM**

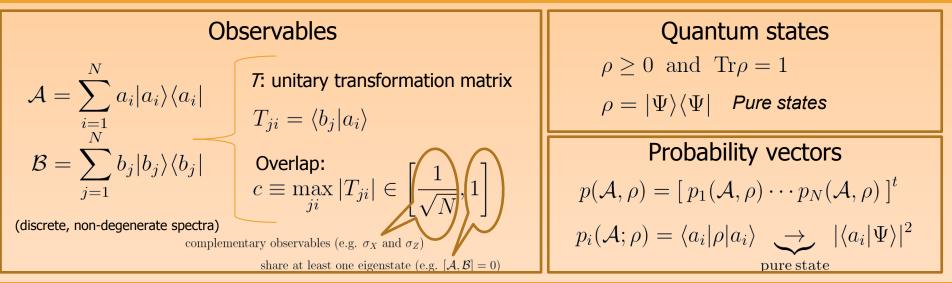
$$\mathcal{U}(\mathcal{A},\mathcal{B},\rho) \geq \mathfrak{B}(\mathcal{A},\mathcal{B})$$

Uncertainty measure associated to the probability distributions of the outcomes

- state-independent
- = 0 only if A and B share at least one eigenstate (commute: particular case)
   > 0 in other case
- >0 in other case

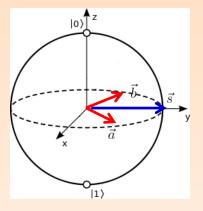
# **UNCERTAINTY RELATIONS**

## **PRELIMINARIES AND NOTATION**



### **HEISENBERG-ROBERTSON RELATION (1929)**

 $V(\mathcal{A}; \rho) V(\mathcal{B}; \rho) \geq \frac{1}{4} |\langle [\mathcal{A}, \mathcal{B}] \rangle_{\rho}|$  $V(\mathcal{A}; \rho) = \sum_{i} (a_{i} - \langle \mathcal{A} \rangle_{\rho})^{2} p_{i}(\mathcal{A}; \rho)$  $\langle \mathcal{A} \rangle_{\rho} = \sum_{i} \overset{i}{a_{i}} p_{i}(\mathcal{A}; \rho)$ 



qubit system 
$$(N = 2)$$
  
 $\rho = \frac{1}{2}(I + \vec{s} \cdot \vec{\sigma}), \ \vec{\sigma} = (\sigma_X, \sigma_Y, \sigma_Z), \ \|\vec{s}\| \le 1$   
 $\mathcal{A} = \vec{a} \cdot \vec{s}, \ \|\vec{a}\| = 1$   $\mathcal{B} = \vec{b} \cdot \vec{s}, \ \|\vec{b}\| = 1$   
HR UR:  
 $\sqrt{1 - (\vec{a} \cdot \vec{s})^2} \sqrt{1 - (\vec{b} \cdot \vec{s})^2} \ge \left| (\vec{a} \times \vec{b}) \cdot \vec{s} \right|$ 

# PART I: ENTROPIC UNCERTAINTY RELATIONS

S. Zozor, G.M. Bosyk, and M. Portesi, *General entropic uncertainty relations for N-level systems,* to be submitted (2013) S. Zozor, G.M. Bosyk, and M. Portesi, *On a generalized entropic uncertainty relation for the qubit,* J. Phys. A: Math. Theo. **46** 465301 (2013)

# **ENTROPY AND UNCERTIANTY RELATIONS**

## **RÉNYI ENTROPY**

1

$$H_{\lambda}[p] = \frac{1}{1-\lambda} \ln \sum_{i=1}^{N} p_{k}^{\lambda}, \quad \lambda \geq 0$$
  
*p*: probability vector  $\{p \in \mathcal{P} = [p_{1} \cdots p_{N}]^{t} \in [0,1]^{N}, \sum_{k} p_{k} = 0$   
 $\Rightarrow$ Shannon entropy:  $H_{1}[p] = -\sum_{i} p_{i} \ln p_{i}$   
 $\Rightarrow$ Colision entropy:  $H_{2}[p] = -\ln \sum_{i}^{i} p_{i}^{2}$   
 $\Rightarrow$ Min entropy:  $H_{\infty}[p] = -\ln \max_{i}^{i} p_{i}$ 

$$p_{k} = \delta_{ki} : H_{\lambda}[p] = 0 \qquad p = \left\lfloor \frac{1}{N} \cdots \frac{1}{N} \right\rfloor^{t} : H_{\lambda}[p] = \ln N$$

$$\diamond H_{\lambda}[p]$$
 is Schur-concave

 $\diamond H_{\lambda}[p]$  is concave in p for  $\lambda \in [0,1]$ 

 $\diamond H_{\lambda}[p]$  for fixed p decreases with  $\lambda$ 

## **GOAL AND PREVIUOS RESULTS**

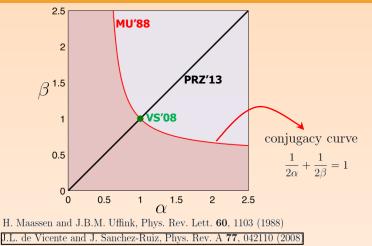
$$H_{\alpha}\left[p(\mathcal{A},\rho)\right] + H_{\beta}\left[p(\mathcal{B},\rho)\right] \ge \overline{\mathfrak{B}}_{\alpha,\beta} > 0, \ (\alpha,\beta) \in \mathbb{R}^{2}_{+}$$

#### T-optimal bound

$$\mathfrak{B}_{\alpha,\beta}(T) = \min_{\substack{\rho \\ \text{Solved only for qubit systems } (N=2), \\ \text{SZ, GMB, MP, JPA 46, 465301 (2013)}} (H_{\alpha}[p(\mathcal{A}, \rho)] + H_{\beta}[p(\mathcal{B}, \rho)]) \geq \overline{\mathfrak{B}}_{\alpha,\beta}$$

#### c-optimal bound

$$\widetilde{\mathfrak{B}}_{\alpha,\beta}(c,N) = \min_{T \in \mathrm{U}(N),\max|T_{ji}|=c} \mathfrak{B}_{\alpha,\beta}(T)$$



Z. Puchała, Ł. Rudnicki, and K. Życzkowski, J. Phys. A: Math. Theor. 46, 272002(2013)

# **METHOD: CONSTRAINED MINIMIZATION**

 $\begin{array}{c} H_{\alpha}\left[p(\mathcal{A},\rho)\right] + H_{\beta}\left[p(\mathcal{B},\rho)\right] \geq \overline{\mathfrak{B}}_{\alpha,\beta}(c) > 0, \ (\alpha,\beta) \in R_{+}^{2} \\ \text{Landau-Pollak inequality (LPI): } \arccos \sqrt{P_{\mathcal{A};\rho}} + \arccos \sqrt{P_{\mathcal{B};\rho}} \geq \arccos c \\ \max \text{maxima probabilities}_{P_{\mathcal{A};\rho}} = \max_{i} \langle a_{i} | \rho | a_{i} \rangle \text{ and } P_{\mathcal{B};\rho} = \max_{j} \langle b_{j} | \rho | a_{j} \rangle \\ \text{Step 1} \\ H_{\lambda,\min}(P) = \min_{p \in \mathcal{P}} H_{\lambda}(p) \text{ s.t. } \max_{k} p_{k} = P \in \left[\frac{1}{N}, 1\right] \\ \text{Equivalent problem:} \\ \overline{\mathfrak{B}}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{B}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{f}_{\alpha,\beta}(c) = \underbrace{\left[ \max_{p \in \mathcal{P}} H_{\alpha,\min}(P_{\mathcal{A}}) + H_{\beta,\min}(P_{\mathcal{A}}) \right]}_{\alpha,\beta} \\ \text{for all } \mathbf{$ 

Equivalent problem:  

$$p = [P \ q_{1} \cdots q_{N-1}]^{t} \text{ and find } q \in \mathbb{R}^{N-1} \text{ that maximizes}$$

$$f_{\lambda}(q) = \frac{\sum_{k} q_{k}^{\lambda} - (1-P)}{\lambda - 1} \quad \text{s.t.} \quad \substack{q \in \mathcal{HC}_{P} \equiv [0; P]^{N-1}}{q \in \mathcal{HP}_{P} \equiv \left\{q : \sum_{k} q_{k} = 1-P\right\}}$$

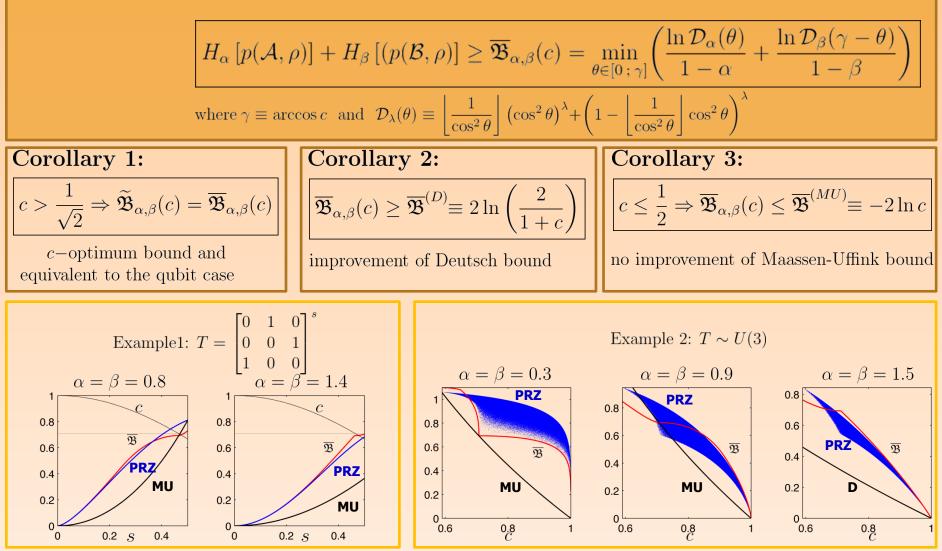
$$\bullet \mathcal{HC}_{P} \cap \mathcal{HP}_{P}: \text{ convex polytope of pure points}$$

$$\left[\underbrace{P \cdots P}_{M \text{ times}} \quad 1 - MP \quad \underbrace{0 \cdots 0}_{N-M-1 \text{ times}}\right]^{t}, \quad M = \left[\frac{1}{P}\right]$$

$$\bullet \int_{0}^{s} \int_{0}$$

# **RESULTS AND EXAMPLES**

### **Proposition:** $\forall (\alpha, \beta) \in \mathbb{R}^2_+$



# PART II: FIDELITY-BASED UNCERTAINTY RELATIONS

G.M. Bosyk, T.M. Osán, P.W Lamberti and M. Portesi, *Geometric formulation of the uncertainty principle*, arXiv:1308.4029v2 [quant-ph] (2013)

# LANDAU-POLLAK INEQUALITY FOR PURE STATES

$$\operatorname{arccos} \left( \begin{array}{c} P_{\mathcal{A};\Psi} + \operatorname{arccos} \left( \begin{array}{c} P_{\mathcal{B};\Psi} \end{array} \right) \geq \operatorname{arccos} c \\ \\ \text{maxima probabilities} \end{array} \right) = \left| \begin{array}{c} P_{\mathcal{A};\Psi} = \max_{i} |\langle a_{i} | \Psi \rangle|^{2} \\ P_{\mathcal{B};\Psi} = \max_{i} |\langle b_{i} | \Psi \rangle|^{2} \end{array} \right|$$

 $\diamond$  if c < 1 then  $P_{\mathcal{A};\Psi}$  and  $P_{\mathcal{B};\Psi}$  cannot both be equal to one

$$\diamond$$
 if  $c = \frac{1}{\sqrt{N}}$  and  $P_{\mathcal{A};\Psi} = 1$  then  $P_{\mathcal{B};\Psi} = \frac{1}{N}$ 

#### Sketch of the proof:

Step 1: consider Wootters metric between state  $|\Psi\rangle$  and  $|a_i\rangle$ ,  $|b_j\rangle$  eigenstates of  $\mathcal{A}$  and  $\mathcal{B}$ 

Wotters metric:  $d_{\rm W}(\varphi, \psi) = \arccos |\langle \varphi | \psi \rangle|$ 

Step 2: apply the triangle inequality  $T_{ij}$  $\operatorname{arccos}(\langle a_i | \Psi \rangle) + \operatorname{arccos}(\langle b_j | \Psi \rangle) \geq \operatorname{arccos}(\langle a_i | b_j \rangle)$ 

Step 3: choose the eigenstates that maximize the probabilities  $\operatorname{arccos} \sqrt{P_{\mathcal{A};\Psi}} + \operatorname{arccos} \sqrt{P_{\mathcal{B};\Psi}} \ge \operatorname{arccos} |\langle a_{i'}|b_{j'}\rangle|$ 

Step 4: use decreasesing property of arccos  $|\langle a_{i'}|b_{j'}\rangle| \ge \arccos\max_{ij} |\langle a_i|b_j\rangle|$ 

H.J. Landau and H.O. Pollak, Bell Syst. Tech. J. **40**, 65 (1961) J. Uffink, Ph. D. thesis, University of Utrecht (1990)

# FIDELITY & PURIFICATION

### FIDELITY

measure of similarity of quantum states

 $F(\rho,\sigma) = \left(\mathrm{Tr}\sqrt{\sqrt{\rho}\sigma\sqrt{\rho}}\right)^2$ 

- normalization  $0 \le F(\rho, \sigma) \le 1$
- identity of indiscernibles  $F(\rho,\sigma) = 1$  iff  $\rho = \sigma$
- symmetry  $F(\rho, \sigma) = F(\sigma, \rho)$
- both pure sates  $F(|\psi\rangle\langle\psi|,|\varphi\rangle\langle\varphi|) = |\langle\psi|\varphi\rangle|^2$
- one pure state  $F(|\psi\rangle\langle\psi|,\sigma)=\langle\psi|\sigma|\psi\rangle$

### PURIFICATION

mixed quantum states  $\rightarrow$  pure quantum states

 $\rho \in B(\mathcal{H}) \to |\Psi\rangle \in \mathcal{H} \otimes \mathcal{H}_{aux} : Tr_{aux}(|\Psi\rangle\langle\Psi|) = \rho$ 

Link between fidelity and purifications

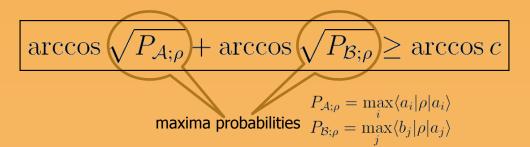
 $F(\rho,\sigma) = \max_{|\varphi\rangle} |\langle \psi | \varphi \rangle|^2$ 

all purifications of  $\sigma$ 

fixed purification  $\rho$ 

## LANDAU-POLLAK INEQUALITY FOR MIXED STATES

#### We prove:



### Sketch of the proof:

Step 1: consider Wootters metric between  $|s\rangle$ ,  $|r\rangle$  and  $|t\rangle$  purifications of  $\sigma$ ,  $\rho$  and  $\tau$ , resp.

Step 2: use triangle inequality  $\arccos |\langle s|r\rangle| + \arccos |\langle t|r\rangle| \ge \arccos |\langle s|t\rangle|$ 

Step 3: choose the purifications that maximize the overlaps with  $\rho$ :  $F(\sigma, \rho) = \max_{|s\rangle} |\langle s|r\rangle|^2 = |\langle \tilde{s}|r\rangle|^2$  and  $F(\tau, \rho) = \max_{|t\rangle} |\langle t|r\rangle|^2 = |\langle \tilde{t}|r\rangle|^2$  $\arccos \sqrt{F(\sigma, \rho)} + \arccos \sqrt{F(\tau, \rho)} \ge \arccos |\langle \tilde{s}|\tilde{t}\rangle|$ 

Step 4: use decreasing property of arccos  $|\langle \tilde{s} | \tilde{t} \rangle| \ge \arccos \sqrt{F(\sigma, \tau)}$ 

#### we proof triangle inequality for the angle!!!

Step 5: use that 
$$\begin{array}{c} \sigma = |a_i\rangle\langle a_i| \\ \tau = |b_j\rangle\langle b_j| \end{array} \xrightarrow{F(\sigma,\rho) = \langle a_i|\rho|a_i\rangle = p_i(\mathcal{A};\rho)} F(\tau,\rho) = \langle b_j|\rho|b_j\rangle = p_j(\mathcal{B};\rho) \end{array} \xrightarrow{\operatorname{arccos}} \operatorname{arccos} \sqrt{p_i(\mathcal{A};\rho)} + \operatorname{arccos} \sqrt{p_j(\mathcal{B};\rho)} \ge \operatorname{arccos} |\langle a_i|b_j\rangle|$$

Step 6 and 7: similar to LPI for pure states, i.e, choose max probabilities and use that arccos decreases

## FIDELITY-BASED UNCERTAINTY RELATIONS

### Metric between quantum states

non-negativy

 $d(\rho, \sigma) \geq 0$  and  $d(\rho, \sigma) = 0$  iff  $\rho = \sigma$ 

symmetry

 $d(\rho, \sigma) = d(\sigma, \rho)$ 

• triangle inequality  $d(\sigma, \rho) + d(\tau, \rho) \ge d(\sigma, \tau)$ 

#### **Fidelity-Based metrics**

Examples:  $d(\rho, \sigma) = f(F(\rho, \sigma))$ • Angle  $d_A(\rho, \sigma) = \arccos \sqrt{F(\rho, \sigma)}$ 

with  $f \downarrow$  and f(x) = 0 iff x = 1 • Bures  $d_{\rm B}(\rho, \sigma) = \sqrt{2 - 2\sqrt{F(\rho, \sigma)}}$ 

• Root-infidelity  $d_{\rm RI}(\rho,\sigma) = \sqrt{1 - F(\rho,\sigma)}$ 

#### Uncertainty measure

$$\mathcal{U}(\mathcal{A};\rho) = f\left(P_{\mathcal{A};\rho}\right)$$

 $\diamond \mathcal{U}(\mathcal{A}; \rho) > 0$  $\diamond \mathcal{U}(\mathcal{A}; \rho)$  is decreasing in terms of  $P_{\mathcal{A}; \rho}$  $\diamond \mathcal{U}(\mathcal{A}; \rho) = 0$  if and only if  $P_{\mathcal{A};\rho} = 1$  $\diamond \arg \max_{P_{\mathcal{A};\rho}} \mathcal{U}(\mathcal{A};\rho) = \frac{1}{N}$ 

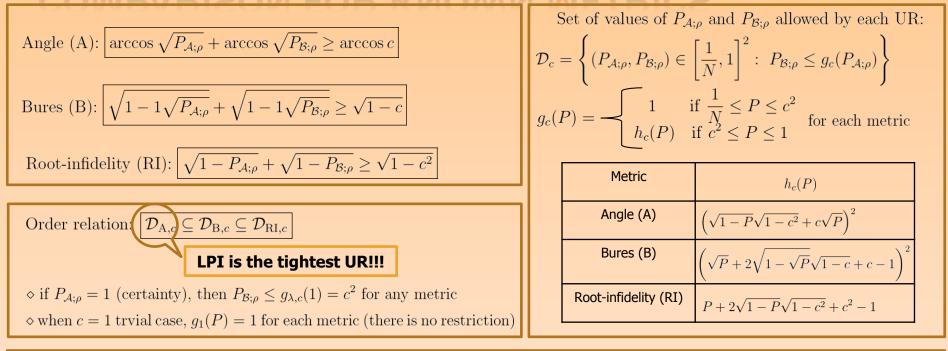
#### **MAIN RESULT**

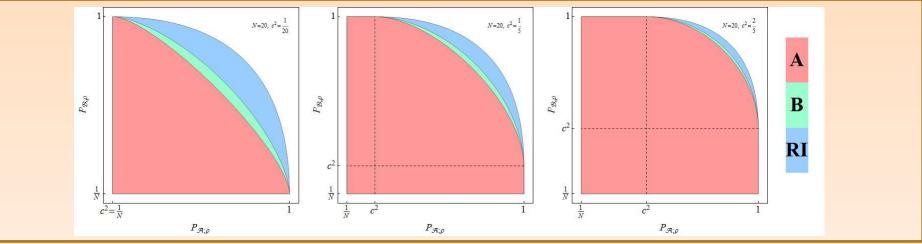
 $\mathcal{A}$  and B observables, and  $\rho$  a denisty operator:  $\mathcal{U}(\mathcal{A};\rho) + \mathcal{U}(\mathcal{B};\rho) \geq f(c^2)$  where  $c = \max_{ij} |\langle b_j | a_i \rangle| \in \left|\frac{1}{\sqrt{N}}, 1\right|$ 

- $\diamond$  family of uncertainty relations
- $\diamond$  universal (state-independent) lower bound
- $\diamond$  for c < 1: uncertainty-sum is strictly greater than zero (non-trivial)

 $\diamond$  for  $c = \frac{1}{\sqrt{N}}$  (complementary observables): certainty on one implies maximum ignorance on the other

## **COMPARISON FOR KNOWN METRICS**







We revisit the formulation of Uncertainty Principle:
 going beyond Heisenberg-like uncertainty relations

□ General Entropic Uncertainty Relations for N-level systems

 $\checkmark$  lower bound for the sum of Rényi entropies of arbitrary entropic indices

 $\checkmark$  c-optimal bound for overlap > 1/ $\sqrt{2}$  (improves all c-dependent bounds,

e.g. Deutsch, MU, etc)

✓ improves PRZ bound in several situations

✓ easy to calculate (solve a one-dimensional minimization problem)
 x open: extension to POVM measurements (LPI to POVM?)

Geometric formulation of Uncertainty Principle

 $\checkmark$  extension of LPI to mixed states

✓ family of fidelity-based uncertainty relations

x conjecture: angle metric gives the tightest uncertainty relation (LPI) within this framework

x open: extension to POVM measurement