A functional Approach for On Line Analytical Processing

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Abstract
This paper describes an approach to On Line Analytical Processing (OLAP), expresed in the declarative programming paradigm.

We define a collection of functions that capture some of the functionality currently provided by multidimensional database product. This is done by defining operations which allow for classifying and reducing relations (tables). Suitably combined, these operations will make possible to carry out the multidimensional analysis of a relational database, and make possible the declarative specification and optimization of multidimensional database queries.

The library works over an abstract model of the relational database calculus as defined by Maier, written in the style of model-oriented formal specification in the functional language Haskell (details can be found in [8]).

Keywords: functional programming, declarative programming, relational data model, on line analytical processing, multidimensional analysis.

Resumen
Este paper describe una aproximación al Procesamiento Analítico On Line (OLAP), desarrollado en el paradigma de la programación declarativa.

Definimos un conjunto de funciones que capturan algunas de las funcionalidades actualmente provistas por las bases de datos multidimensionales existentes. Esto es realizado definiendo operaciones que permiten clasificar y reducir relaciones (tablas). Dichas operaciones, convenientemente combinadas, permitirán llevar a cabo el análisis multidimensional de una base de datos relacional, junto con la especificación declarativa y optimización de las correspondientes consultas multidimensionales sobre dicha base.

La biblioteca provista trabaja sobre un modelo abstracto del cálculo de base de datos relacionales tal como ha sido definido por Maier,escrito en un estilo de especificación formal orientada al modelo en el lenguaje funcional Haskell (más detalles pueden encontrarse en [8])


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1 INTRODUCTION

Codd proposed the concept of On-Line Analytical Processing (OLAP) for rendering enterprise data in multidimensional perspectives, performing on-line analysis of data using mathematical formulas or more sophisticated statistical analyses, and consolidating and summarizing data [3], [4].

OLAP call for sophisticated on-line analysis, something for which the traditional relational model [2] offers little support. Several vendors have already developed OLAP products, but many of these suffer from the following limitations: they do not support a comprehensive “query” language similar to SQL; viewing data in multi-dimensional perspectives involves treating certain attributes as dimensional parameters and the remaining ones as measures, and then analyzing them as a “function” of the parameters; and, finally, unlike for the relational model, there is no precise, commonly agreed, conceptual model for OLAP or the so-called multidimensional databases (MDD) (see [5], [1], [6]).

We present a comprehensive, simple conceptual model for OLAP that treat dimensions and measures symmetrically.

The structure of the paper as follows: Section 2 introduces elementary concepts and terminology which are used throughout the paper. Section 3 sketches a formal model for database relational data. Section 4 contains a very brief comparison between relational and multidimensional tables. Section 5 presents the functionality necessary for OLAP-based applications. The last section present some conclusions and future work.

2 CATEGORICAL SUPPORT

Categories. A category consists of a collection of objects and a collection of arrows. Each arrow \( f :: a \rightarrow b \) has a source object \( a \) and a target object \( b \). Two arrows \( f \) and \( g \) can be composed to form a new arrow \( g \cdot f \), if \( f \) has the same target object as the source object of \( g \). This composition operation is associative. Furthermore, for each object \( a \) there is a so-called identity arrow \( id_a :: a \rightarrow a \), which is the unit of composition.

Our base category is called \( Types \) and has types as objects and functions as arrows. Arrow composition is function composition (\( . \)) and the identity arrows are represented by the polymorphic function \( id \).

Functors. Functors are structure-preserving mappings between categories. Polymorphic datatypes are functors from \( Types \) to \( Types \). In Haskell, functors can be defined by a type constructor \( f \) of kind \( * \rightarrow * \), mapping objects to objects, together with a higher-order function \( fmap \), mapping arrows to arrows. This is provided as a constructor \( class \) in the Haskell Prelude (the standard file of primitive functions) as follows:

\[
\text{class Functor f where}
\text{fmap :: (}a \rightarrow b\text{) } \rightarrow \text{ (}f \text{ a } \rightarrow f \text{ b)}
\]

The arrow action of a functor must preserve identity arrows and distribute over arrow composition. For functors from \( Types \) to \( Types \), this means that the following equations must hold:

\(fmap \ id = \ id\)
\[ \text{fmap} (f \cdot g) = (\text{fmap} f) \cdot (\text{fmap} g) \]

**Bifunctors.** The product category \( \text{Types} \times \text{Types} \) consists of pairs of types and pairs of functions. We can define functors from \( \text{Types} \times \text{Types} \) to the base category \( \text{Types} \) in Haskell. These functors are called *bifunctors*. A (curried) bifunctor in Haskell is a type constructor of kind \( * \to * \to * \), together with a function \( \text{bmap} \). The following constructor class \( \text{Bifunctor} \) was made available:

```haskell
class Bifunctor f where
  bmap :: (a -> c) -> (b -> d) -> (f a b -> f c d)
```

**Products.** Categorical products are provided in Haskell by the type constructor for pairs \( (a, b) \) (usually written as Cartesian product \( a \times b \) in mathematics) and projections \( \text{fst} \) and \( \text{snd} \) (resp. \( \pi_1 \) and \( \pi_2 \) in standard mathematical notation). Type constructor \( (, \) is extended to a bifunctor in the obvious way:

```haskell
instance BiFunctor (,) where
  bmap f g = f \times g
```

where

\[
(\times) :: (a \to b) \to (c \to d) \to (a, c) \to (b, d)
\]

\[
(f \times g) = \text{split} (f \cdot \text{fst}) (g \cdot \text{snd})
\]

and combinator \( \text{split} :: (a \to b) \to (a \to c) \to a \to (b, c) \) behaves as follows: \( \text{split} f g x = (f x, g x) \).

**Sums.** Categorical sums are defined in the Haskell Prelude by means of type constructor

```
data Either a b = Left a | Right b
```

together with a function \( \text{either} :: (a \to b) \to (c \to b) \to \text{Either} a c \to b \) satisfying the following equations:

\[
(\text{either} f g) \cdot \text{Left} = f
\]

\[
(\text{either} f g) \cdot \text{Right} = g
\]

Type constructor \( \text{Either} \) is extended to a bifunctor by providing the following instance of \( \text{bmap} \):

```
(+:: (a -> b) -> (c -> d) -> Either a c -> Either b d
(f + g) (Left a) = Left (f a)
(f + g) (Right b) = Right (g b)
```

```
instance BiFunctor Either where
  bmap f g = f + g
```

The popular notations \( \langle f, g \rangle \), \( [f, g] \) and \( F f \) (where \( F \) is a functor) will be adopted interchangeably with \( \text{split} f g \), \( \text{either} f g \) and \( \text{fmap} f \), respectively.
Invertible arrows. An arrow \( f :: b \to a \) is said to be right-invertible (vulg. surjective) if there exists some \( g :: a \to b \) such that \( f \cdot g = \text{id}_a \). Dually, \( g \) is said to be left-invertible (vulg. injective) if there exists some \( f \) such that the same fact holds. Then type \( b \) is said to “represent” type \( a \) and we draw:

\[
\begin{array}{c}
a \downarrow g \\
\downarrow f \\
\le \\
b
\end{array}
\]

where \( g \) and \( f \) are called resp. the representation and abstraction functions. An isomorphism \( f :: b \to a \) is an arrow which has both a right-inverse \( g \) and a left-inverse \( h \) — a bijection in set theory terminology. It is easy to show that \( g = h = f^{-1} \). Type \( a \) is said to be isomorphic to \( b \) and one writes \( a \cong b \).

Isomorphisms are very important functions because they convert data from one “format” to another format losing information. These formats contain the same “amount” of information, although the same datum adopts a different “shape” in each of them. Many isomorphisms useful in data manipulation can be defined [10], for instance function \( \text{swap} :: (a, b) \to (b, a) \) which is defined by \( \text{swap} = \langle \pi_2, \pi_1 \rangle \) and establishes the commutative property of product, \( a \times b \cong b \times a \).

3 MODELING RELATIONAL DATA

Collective datatypes. Our model of relational data will be based on several families of abstractions, including collective datatypes such as finite powersets \((P_a)\) and finite partial mappings \((a \leftarrow b)\). These are modeled as Haskell polymorphic algebraic types (that is, algebraic type definitions with type variables) based on finite lists \([a]\), see \( \text{Set} a \) and \( P\text{fun} a b \) in Table 1, respectively. Both abstractions contain an equality relation and an ordering relation. The latter instantiates to set inclusion \((\subseteq)\) and partial function definedness, respectively.

The finite sets model assumes invariant \( \phi (\text{Set} l) \overset{\text{def}}{=} \text{length} l = \text{card}(\text{elems} l) \), where \( \text{length} \) is Haskell standard and \( \text{card}(\text{inal}) \) and \( \text{elem}(\text{ent})s \) have the usual set-theoretical meaning. Partial mappings require an extra invariant ensuring a functional dependence on sets of pairs \(^1\):

\[
\text{fdp} \overset{\text{def}}{=} (\subseteq \{1\}) \cdot \text{rng} \cdot (\text{id} \to \text{card}) \cdot \text{collect}
\]

Table 1 summarizes the Haskell modules defined for these datatypes.

Relational Database Model. An \( n \)-ary relation in mathematics is a subset of a finite \( n \)-ary product \( A_1 \times \ldots \times A_n \), which is inhabited by \( n \)-ary vectors \((a_1, \ldots, a_n)\). Each entry \( a_i \) in vector \( t = (a_1, \ldots, a_n) \) is accessed by its position’s projection \( \pi_i : A_1 \times \ldots \times A_n \to A_i \). This, however, is not expressive enough to model relational data as this is understood in database theory [7]. Two ingredients must be added, whereby vectors give place to tuples: attribute names and NULL values. Concerning the former, one starts by rendering vectorial indices explicit, in the sense of writing \( e.g. t \ i \) instead of \( \pi_i \ t \). This implies merging all datatypes \( A_i \) to \( A_n \) into a single coproduct type \( A = \sum_{i=1}^n A_i \) and then represent the \( n \)-ary product as:

\[
A_1 \times \ldots \times A_n \leq (\sum_{i=1}^n A_i)^n
\]

\(^1\text{collect} :: P(a \times b) \to (a \to P b) \) converts a relation into a set-valued partial function and \( \text{rng} :: (a \to b) \to P b \) is the usual range function.
| Datatypes: | data Set a = Set [a] | data P fun a b = Map[(a, b)] |
| Constructors: | emptyS, sings, puts, prods llos | bottom, singpf, putpf collect |
| Deletions: | gets | getpf |
| Observers: | ins, nins, incls, card allS | compatible, incompatible allP f |
| Filters: | filterS | |
| Operations: | inters, unions, diffS, plus, pf zip flatr, flatl, slstr, srstr, sextl, sextr zipS, zipWithAllS | plus, pf inv, restn, restp pf zip, pf zipWith |
| Folds: | foldS | foldPF |
| Functor: | fmapS | bmapP f |
| Bifunctor: | | |
| Others: | the, stol, elems, card unzipS | dom, rng, aplpf tnest, discollect, mkr, bpfTrue, bpfFalse pf unzip |

Table 1: Finite sets and partial functions: datatypes and functions implemented.

under representation function $^2 r \ (a_j)_{j=1..n} \overset{\text{def}}{=} \lambda j. (i_j \ a_j)$ which entails invariant

$$\phi \ t \overset{\text{def}}{=} \forall j = 1, \ldots, n, t \ j = i_j \ x : x \in A_j$$

Note that $j = 1, \ldots, n$ can be written $j \in \pi$, where $\pi = \{1, \ldots, n\}$ is the initial segment of the natural numbers induced by $n$. Set $\pi$ is regarded as the attribute name-space of the model $^3$.

As a second step in the extension of vectors to tuples, we consider the fact that some attributes may not be present in a particular tuple, that is, NULL values are allowed $^4$:

$$(\sum_{i \in \pi} A_i + 1)^n$$

which finally leads to tuples as inhabitants of

$$\text{T} \! \! \text{uple} \ = \ (\pi \to \sum_{i \in \pi} A_i)$$

thanks to isomorphism $A \rightarrow B \cong (B + 1)^A$ $^9$. This models tuples of arbitrary arity (up to $n$ attributes), including the empty tuple. For notation economy, for every $X \subseteq \pi$, we will write $T \! \! \text{uple}_X$ as a shorthand for $X \rightarrow \sum_{i \in X} A_i$.

$T \! \! \text{uple}$ is the basis for the Haskell model of database relations presented in Table 2. Relations (Relation) are sets of tuples sharing a common attribute schema (SchemaR). A rather complex invariant ensuring that tuples are well and consistently typed is required, which is omitted here for economy of presentation. This and other details of this model can be found in $^8$.

$^2$Injections $i_j$ are associated to the $n$-ary coproduct. Left and Right in Haskell correspond to $i_1$ and $i_2$, respectively.

$^3$The fact that this can be replaced by any isomorphic collection of attribute names of cardinality $n$ has little impact in the modelling, so we stick to $\pi$.

$^4$Think of 1 as the singleton type \{NULL\}.
4 RELATIONAL VERSUS MULTIDIMENSIONAL TABLES

The fundamental data structure of a multidimensional database is what we call an \textit{n-dimensional table}. Let us start by giving some intuition behind the concept. We wish to be able to see values of certain attributes as “functions” of others, in whichever way suits us, exploiting possibilities of multidimensional rendering. Drawing on the terminology of statistical databases [11], we can classify the attribute set associated with the scheme of a table into two kinds: \textit{parameters} and \textit{measures}. There is no a priory distinction between parameters and measures, so that any attribute can play either role\footnote{Needless to say, the data type of a measure attribute must have some kind of metrics or algebra associated with it.}. An example of a two-dimensional table is given in Table 3 (adapted from [5]).

\begin{table}[h]
\centering
\begin{tabular}{|c|c|}
\hline
\textbf{SALES} & \textbf{TIME} \\
\hline
\textbf{Year} & 1996 & 1997 \\
\hline
\textbf{Month} & Jan & Feb ... & Jan & Feb ... \\
\hline
\textbf{Part} & \textbf{City} & \textbf{(Cost, Sale)} & \textbf{(Cost, Sale)} \\
\hline
PC & Mendoza & (5, 6) & (5, 7) ... & (4, 6) & (4, 8) ... \\
\hline
Córdoba & (5, 7) & (5, 8) ... & (4, 8) & (4, 9) ... \\
\hline
... & ... & ... & ... & ... \\
\hline
Inkjet & Mendoza & (7, 8) & (7, 9) ... & (6, 9) & (6, 8) ... \\
\hline
Bs. As. & (6, 9) & (6, 9) ... & (5, 8) & (5, 9) ... \\
\hline
... & ... & ... & ... & ... \\
\hline
\end{tabular}
\caption{\textit{SALES} — a sample two dimensional table with dimensions \textit{Category} and \textit{Time}. The associated parameter sets are \{\textit{Part}, \textit{City}\} and \{\textit{Year}, \textit{Month}\}, respectively. The measure attributes are \textit{Cost} and \textit{Sale}.}
\end{table}

We want to work with the relational model we have defined in the previous section. A natural way to achieve this is to regard the multidimensionality of tables as an inherently \textit{structural} feature, which is most significant when the table is rendered to the user. The actual \textit{contents} of a table are essentially orthogonal to the associated structure, i.e., the distribution of attributes over dimensions and measures. Separating both features leads to a \textit{relational} view of a table. For instance, the entry in the first (i.e., top left-most) “cell” in Table 3 containing the entry \((5, 6)\) corresponds to the tuple...
over the scheme

\{Part, City, Year, Month, Cost, Sale\} \hspace{1cm} (2)

in a relational view of table \textit{SALES}.

To provide for OLAP, we need to define operations concerned with the following kinds of functionality:

- \textbf{Classification}: Ability to classify or group data sets in a manner appropriate for subsequent summarization.

- \textbf{Reduction/Consolidation}: Generalization of the aggregate operators in standard SQL. In general, reduction maps multi-sets of values of a numeric type to a single, “consolidated”, value.

Classification is a generalization of the familiar SQL \textbf{group by} operator. The following example presents a typical query involving classification.

\textbf{Example 4.1} Consider the relation \textit{RSALES} with scheme (2) mentioned before. A typical query \textit{would be}: “find, for each part, the total amount of annual sales”. \textit{Even though this query involves aggregation, notice that it also involves classifying the data into various groups according to certain criteria, before aggregation is applied. Concretely, the above query involves classification by attributes \textit{Part} and \textit{Year}.}

\section{OLAP-BASED APPLICATIONS FRAME}

\subsection{Classification}

In our model, relations are sets of tuples (tables) with a scheme, while tuples are finite partial functions. First, we define a function for partial function decomposition (or “tuple classification”). Then we extend the notion of classification, applying it in the context of tables.

\textbf{Partial function decomposition (Classification on Tuples)} Let \( t \) be a tuple \(( t \in \text{Attribute} \to \text{Value})\), and let \( X = \{A_1, ..., A_k\} \) be an arbitrary subset of \( \text{dom}(t) \). A \textbf{classification} over \( X \) of tuple \( t \), is the pair of tuples defined by \( \text{tnest} \ X \), where \( \text{tnest} \) is polymorphic function

\begin{align*}
\text{tnest} & : \mathcal{P}A \to (A \to B) \to ((A \to B) \times (A \to B)) \\
tnest \ s & \overset{\text{def}}{=} \langle | s, \setminus s \rangle
\end{align*}

The idea of this function is to decompose a partial map into a pair of maps of the same type

\begin{align*}
A \to B \overset{\text{tnest} \ s}{\leq} (A \to B) \times (A \to B)
\end{align*}

which, together, rebuild the original map. In Haskell:

\begin{verbatim}
  tnest :: Eq a => Set a -> Pfun a b -> (Pfun a b,Pfun a b)
  tnest s f = (s <: f,s <-: f)
\end{verbatim}
Tabular Decomposition (Classification on Tables) Let $t$ be a set of tuples (typed $\mathcal{P}(\text{Attribute} \rightarrow \text{Value})$), and let $X = \{A_1, ..., A_k\}$ be an arbitrary set of attributes of $t$. A classification over $X$, of table $t$, is given by $t\text{collect } X \ t$, where $t\text{collect}$ is polymorphic function

\[
t\text{collect } s \overset{\text{def}}{=} \text{collect} \cdot \mathcal{P}(t\text{nest } s)
\]

that is,

\[
t\text{collect } :: (\text{Eq } a, \text{Eq } b) \\
\rightarrow \text{Set } a \\
\rightarrow \text{Set } (\text{Pfun } a b) \\
\rightarrow \text{Pfun } (\text{Pfun } a b) \ (\text{Set } (\text{Pfun } a b))
\]

t\text{collect } s \ t = \text{collect} \ (\text{nmap } (t\text{nest } s) \ t)

in Haskell.

Classification essentially maps tuples of a relation to different groups (necessarily disjoint). Intuitively, we can think of the attributes in the first argument of $t\text{collect}$ as corresponding to the “group id”.

Example 5.1 The classification part of the query of Example 4.1 can be expressed as follows: $t\text{collect} \ \{\text{“Part”}, \text{“Year”}\}$ (tuples RSALES). Table 4 illustrates the result of this operation in concrete Haskell syntax.

5.2 Reduction/Consolidation

Next, we consider reduction/consolidation, which includes not only applications of functions such as $\text{max}$, $\text{min}$, $\text{avg}$, $\text{sum}$, $\text{count}$ to multi-sets of values defined by groups of tuples, but also statistical functions such as $\text{variance}$ and $\text{mode}$, and business calculations such as $\text{proportions}$ and $\text{quarterlies}$.

In our model, reduction functions map sets of tuples of values to individual values. We first define some necessary auxiliary functions.

5.2.1 Relational Reduction

Function

\[
t\text{tot}2 \ : \ (A \times B \rightarrow B) \rightarrow B \rightarrow \mathcal{P}(C \times A) \rightarrow B
\]

\[
t\text{tot}2 \ f \ u \overset{\text{def}}{=} \{[u, f \circ (\pi_2 \times \text{id})]\}
\]

reduces a binary relation on the second projection according to a reduction structure $A \times B + 1 \xrightarrow{(f,u)} B$ which, in most cases, is a monoid algebra. In Haskell:

\[
t\text{tot}2 :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow \text{Set } (c, a) \rightarrow b
\]
\[
t\text{tot}2 \ f = \text{foldS } (\text{curry } (\text{uncurry } f \ . \ (\text{p2 } (\text{op} \rightarrow \text{id}))))
\]
Map[ ( Map ["Part", "PC"], "Year", "1996") , 
Set[ Map ["City", "Mendoza"], ("Month", "Jan"), ("Cost", 5), ("Sale", 6)], 
Map ["City", "Mendoza"], ("Month", "Feb"), ("Cost", 5), ("Sale", 7)], 
.................................................................................. 
Map ["City", "Cordoba"], ("Month", "Jan"), ("Cost", 5), ("Sale", 7)], 
.................................................................................. 
( Map ["Part", "PC"], "Year", "1997") , 
Set[ Map ["City", "Mendoza"], ("Month", "Jan"), ("Cost", 4), ("Sale", 6)], 
Map ["City", "Mendoza"], ("Month", "Feb"), ("Cost", 4), ("Sale", 8)], 
.................................................................................. 
Map ["City", "Cordoba"], ("Month", "Jan"), ("Cost", 4), ("Sale", 8)], 
.................................................................................. 
( Map ["Part", "Inkjet"], "Year", "1996") , 
Set[ Map ["City", "Mendoza"], ("Month", "Jan"), ("Cost", 7), ("Sales", 8)], 
Map ["City", "Mendoza"], ("Month", "Feb"), ("Cost", 7), ("Sale", 9)], 
.................................................................................. 
Map ["City", "Buenos Aires"], ("Month", "Jan"), ("Cost", 6), ("Sale", 9)], 
.................................................................................. 
( Map ["Part", "Inkjet"], "Year", "1997") , 
Set[ Map ["City", "Mendoza"], ("Month", "Jan"), ("Cost", 6), ("Sales", 9)], 
Map ["City", "Mendoza"], ("Month", "Feb"), ("Cost", 6), ("Sale", 8)], 
.................................................................................. 
Map ["City", "Buenos Aires"], ("Month", "Jan"), ("Cost", 5), ("Sale", 8)], 
.................................................................................. ]

Table 4: Output of the expression \( tcollect \{ "Part", "Year" \} \{ tuples RSALES \} \).

5.2.2 Partial Function Application with a Default Value

Let \( apl \) be the isomorphism

\[
A \rightarrow B \xrightarrow{apl} (B + 1)^A \xrightarrow{apl^{-1}} A
\]

in

\[
get : B \rightarrow A \rightarrow (A \rightarrow B) \rightarrow B
\]

\[
get u a f \overset{\text{def}}{=} [id, u] \cdot (apl f a)
\]

In Haskell:

\[
get :: Eq a \Rightarrow b \rightarrow a \rightarrow Pfun a b \rightarrow b
\]

\[
get u a f = aux (aplpf' f a)
\]

where \( aux \ (Ok b) = b \)

\[
aux \ (Err s) = u
\]
5.2.3 Tabular Reduction

Finally, function
\[
ttot : B \rightarrow ((A \times A) \rightarrow A) \rightarrow A \rightarrow \mathcal{P}(B \rightarrow A) \rightarrow (B \rightarrow A)
\]

\[
ttot b f u s \triangleq \{ b \mapsto \text{tot2 } f u (\mathcal{P}(g u b)) s \}
\]

performs tabular reduction, where \( g u b \triangleq (id \times (\text{get } u b)) \cdot \text{swap} \cdot (\text{tnest } \{ b \}) \). Argument \( b \) specifies the measure attribute over which reduction will take place while arguments \( f \) and \( u \) provide the required reduction algebra. The output is packaged into a one-attribute tuple mapping the measure attribute name to the final result.

The corresponding Haskell code follows the above definition very closely:

\[
ttot :: (\text{Eq } a, \text{Eq } b) \\
\quad -> b \\
\quad -> (a -> a -> a) \\
\quad -> a \\
\quad -> \text{Set } (\text{Pfun } b a) \\
\quad -> \text{Pfun } b a \\
\]

\[
ttot b f u s = \text{Map } [ b |-> (\text{tot2 } f u (g u b s)) ]
\]

where \( g u b = \text{nmap } ((\text{id } \bowtie (\text{get } u b)) \cdot \text{swap} \cdot \text{tnest } (\text{sings } b)) \)

**Example 5.2** Consider again the query of Example 4.1. We illustrate in this example how, from the classified set of tuples computed in Example 5.1, it is possible to obtain the final answer to the query.

Let \( f\text{lass} \) be the mapping arising from the classification step (\( f\text{lass} \) type is \((\text{Attr } \rightarrow \text{Value}) \rightarrow \mathcal{P}(\text{Attr } \rightarrow \text{Value})) \) computed in Example 5.1. We can use \( ttot \) to summarize over the \textit{Sales} attribute, with a particular binary operation (monoid \((+, 0)\) in this example), in the range of \( f\text{lass} \). The last step is to transform the resulting structure in a table (set of tuples). Diagram (5) depicts the required computations.

\[
\mathcal{P}(\text{Attr } \rightarrow \text{Value}) \\
\quad \downarrow \text{tcollect } \{ \text{Part,Year} \} \\
(\text{Attr } \rightarrow \text{Value}) \rightarrow \mathcal{P}(\text{Attr } \rightarrow \text{Value}) \\
\quad \downarrow \text{id } \rightarrow (ttot \text{ Sales } (+, 0)) \\
(\text{Attr } \rightarrow \text{Value}) \rightarrow (\text{Attr } \rightarrow \text{Value}) \\
\quad \downarrow \text{mkr} \\
\mathcal{P}( \text{(Attr } \rightarrow \text{Value}) \times (\text{Attr } \rightarrow \text{Value})) \\
\quad \downarrow \mathcal{P}(\text{uncurry plus}) \\
\mathcal{P}(\text{Attr } \rightarrow \text{Value})
\]

Altogether, we have evaluated the expression

\[
(\mathcal{P}\text{plus} ) \cdot \text{mkr} \cdot (\text{id } \rightarrow (ttot \text{ Sale } (+, 0))) \cdot (\text{tcollect } \{ \text{Part,Year} \})
\]
5.3 Multidimensional Analysis

Next we define a “multidimensional analysis” function that generalizes the algebraic structure of (6) above:

\[
mda \ s \ a \ f \ u \ \overset{\text{def}}{=} (P_{\text{plus}}) \cdot mkr \cdot (B \rightarrow t\text{tot} \ a \ f \ u) \cdot (t\text{collect} \ s)
\]

In the context of our relational model in Haskell, we provide the \( mda \) function defined over the \textit{Relation} data type, as follows:

\[
mdaR :: \text{Set IdAttr} \\
\rightarrow \text{IdAttr} \\
\rightarrow (\text{Value} \rightarrow \text{Value} \rightarrow \text{Value}) \\
\rightarrow \text{Value} \\
\rightarrow \text{Relation} \\
\rightarrow \text{Relation}
\]

\[
mdaR \ s \ a \ f \ u \ r = \\
\text{Rel} \ ((\text{unions} \ s \ (\text{sings} \ a)) <:(\text{schema} \ r)) \\
(\text{nmap} \ (\text{uncurry} \ \text{plus})(\text{mkr} \ y)) \\
\text{where} \ y = (\text{id} \ *\rightarrow \ (\text{ttot} \ a \ f \ u)) \ x \\
x = \text{tcollect} \ s \ (\text{tuples} \ r)
\]

Table 5 illustrates the application of the “multidimensional analysis” operation \( mdaR \) to our running example. Operation \( mdaR \) produces a relation with scheme \textit{Part}, \textit{Year}, \textit{Sale} which is depicted two-dimensionally.

<table>
<thead>
<tr>
<th>PartYearSales</th>
<th>Year</th>
</tr>
</thead>
<tbody>
<tr>
<td>1996</td>
<td>1997</td>
</tr>
<tr>
<td>\textbf{Part}</td>
<td>\textbf{PC}</td>
</tr>
<tr>
<td>Inkjet</td>
<td>298</td>
</tr>
</tbody>
</table>

Table 5: Output of the expression \( mdaR \ (\text{Set \{"Part","Year"\}}) \ ) \ "Sale" \ f\text{add} \ f0 \ RSALES \) applied to the input relation \textit{RSALES} of Example 4.1.

6 CONCLUDING REMARKS

The research carried out in this paper belongs to the intersection of formal methods with relational database theory.

In this paper we have presented a collection of functions in Haskell language that capture some of the functionality currently provided by multidimensional database products. We worked under the assumption that relational systems can model N-dimensional data as a relation with N-attribute domains.

Function \( mda \), which creates a table with an aggregated value indexed by a set of attributes, operates on relations and produces relations. It could be composed with the basic operators of relational algebra to build other \textit{OLAP} operators in order to provide constructs such histograms, cross-tabulations, subtotals, roll-up and drill-down. For instance, \( mda \) could be used to compute the following table (roll up using totals report):

The rightmost column corresponds to the output of the expression:
Table 6: Sales Roll Up by Part by Year (applied to the input relation RSALES of Example 4.1.

<table>
<thead>
<tr>
<th>Part</th>
<th>Year</th>
<th>Total</th>
</tr>
</thead>
<tbody>
<tr>
<td>PC</td>
<td>1996</td>
<td>320</td>
</tr>
<tr>
<td></td>
<td>1997</td>
<td>455</td>
</tr>
<tr>
<td></td>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td></td>
<td>1256</td>
<td></td>
</tr>
<tr>
<td>Inkjet</td>
<td></td>
<td>298</td>
</tr>
<tr>
<td></td>
<td>1996</td>
<td>450</td>
</tr>
<tr>
<td></td>
<td>1997</td>
<td>...</td>
</tr>
<tr>
<td></td>
<td>987</td>
<td></td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>1788</td>
</tr>
<tr>
<td></td>
<td>1450</td>
<td></td>
</tr>
<tr>
<td></td>
<td>...</td>
<td>—</td>
</tr>
</tbody>
</table>

mdaR (Set["Part"]) "Sale" fadd f0 RSALES

while the bottom row could be computed using:

mdaR (Set["Year"]) "Sales" valadd val0 r1olap

Composition (and others categorical combinators) provides a powerful tool to compose operators and allows for complex multidimensional queries to be built.

It should be stressed that the operations defined do not intend to address the issue of restructuring information from the perspective of the dimensionality of the data.

The next step, which goes beyond the scope of the present work, should include the development of a complete algebra and a rigorous calculus based upon the algebraic operators of the model.

REFERENCES


